

6.1100

Lecture 1: Introduction

Staff

- Lecturer
 - Prof. Martin Rinard rinard@mit.edu 258-6922 32-G828
 - Course Secretary
 - Mary McDavitt mmcdavit@csail.mit.edu 253-9620 32-G785
 - Teaching Assistants
 - Krit Boonsiriseth (talkon@mit.edu)
 - Pleng Chomphoochan (tcpc@mit.edu)
 - Youran (Yoland) Gao (youran@mit.edu)
 - Tarushii Goel (tarushii@mit.edu)
 - Web Site
 - <https://6110-sp24.github.io/syllabus>

Reference Textbooks

- *Modern Compiler Implementation in Java (Tiger book)*
A.W. Appel
Cambridge University Press, 1998
ISBN 0-52158-388-8

A textbook tutorial on compiler implementation, including techniques for many language features
- *Advanced Compiler Design and Implementation (Whale book)*
Steven Muchnick
Morgan Kaufman Publishers, 1997
ISBN 1-55860-320-4

Essentially a recipe book of optimizations; very complete and suited for industrial practitioners and researchers.
- *Compilers: Principles, Techniques and Tools (Dragon book)*
Aho, Lam, Sethi and Ullman
Addison-Wesley, 2006
ISBN 0321486811

The classic compilers textbook, although its front-end emphasis reflects its age. New edition has more optimization material.
- *Engineering a Compiler (Ark book)*
Keith D. Cooper, Linda Torczon
Morgan Kaufman Publishers, 2003
ISBN 1-55860-698-X

A modern classroom textbook, with increased emphasis on the back-end and implementation techniques.
- *Optimizing Compilers for Modern Architectures*
Randy Allen and Ken Kennedy
Morgan Kaufman Publishers, 2001
ISBN 1-55860-286-0

A modern textbook that focuses on optimizations including parallelization and memory hierarchy optimization

The Project: The Five Segments

- ① Lexical and Syntax Analysis
- ② Semantic Analysis
- ③ Code Generation
- ④ Dataflow Analysis
- ⑤ Optimizations

Each Segment...

- Segment Start
 - Project Description
- Lectures
- Project Time – No Class
 - (Design Document)
 - (Project Checkpoint)
- Project Due

Project Groups

- 1st project is an individual project
- Projects 2 to 5 are group projects
- Each group consists of 3 to 4 students
- Projects are designed to produce a compiler by the end of class

Project Collaboration Policy

- Talk about anything you want with anybody
- Write all the code yourself (with LLMs)
- Check with TAs before using specialized libraries designed to support compiler construction
- ChatGPT/copilot/LLMs
 - We encourage you to use LLMs as much as you like
 - You can use/copy/modify anything you get from LLM
 - It should come from your interactions
 - We will ask you to talk about LLM contributions to assignments you turn in
- See the website for specifics

Quizzes

- Two Quizzes
- Unless we have another remote semester, in which case we may do problem sets or some combination instead

More Course Stuff

- Blank page project – all the rope you want!
- Challenging project
- You are on your own!

Why Study Compilers?

- Compilers enable programming at a high level language instead of machine instructions.
 - Malleability, Portability, Modularity, Simplicity, Programmer Productivity
 - Also Efficiency and Performance
- Indispensible programmer productivity tool
- One of most complex software systems to build

What a Compiler Does

- Input: High-level programming language
- Output: Low-level assembly instructions
- Compiler does the translation:
 - Read and understand the program
 - Precisely determine what actions it requires
 - Figure-out how to faithfully carry out those actions
 - Instruct the computer to carry out those actions

Input to the Compiler

- Standard imperative language (Java, C, C++)
 - State
 - Variables,
 - Structures,
 - Arrays
 - Computation
 - Expressions (arithmetic, logical, etc.)
 - Assignment statements
 - Control flow (conditionals, loops)
 - Procedures

Output of the Compiler

- State
 - Registers
 - Memory with Flat Address Space
- Machine code – load/store architecture
 - Load, store instructions
 - Arithmetic, logical operations on registers
 - Branch instructions

Compilers Optimize Programs for...

- Performance/Speed
- Code Size
- Power Consumption
- Fast/Efficient Compilation
- Security/Reliability
- Debugging

Example (input program)

```
int sumcalc(int a, int b, int N)
{
    int i, x, y;
    x = 0;
    y = 0;
    for(i = 0; i <= N; i++) {
        x = x + (4*a/b)*i + (i+1)*(i+1);
        x = x + b*y;
    }
    return x;
}
```

Unoptimized Code

```
pushq %rbp
movq %rsp, %rbp
movl %edi, -4(%rbp)
movl %esi, -8(%rbp)
movl %edx, -12(%rbp)
movl $0, -20(%rbp)
movl $0, -24(%rbp)
movl $0, -16(%rbp)
.L2: movl -16(%rbp), %eax
cmpl -12(%rbp), %eax
jg .L3
movl -4(%rbp), %eax
leal 0(%rax, 4), %edx
leaq -8(%rbp), %rax
movq %rax, -40(%rbp)
movl %edx, %eax
movq -40(%rbp), %rcx
cltd
idivl (%rcx)
movl %eax, -28(%rbp)
movl -28(%rbp), %edx
imull -16(%rbp), %edx
movl -16(%rbp), %eax
incl %eax
imull %eax, %eax
addl %eax, %edx
leaq -20(%rbp), %rax
addl %edx, (%rax)
movl -8(%rbp), %eax
movl %eax, %edx
imull -24(%rbp), %edx
leaq -20(%rbp), %rax
addl %edx, (%rax)
leaq -16(%rbp), %rax
incl (%rax)
jmp .L2
.L3: movl -20(%rbp), %eax
leave
ret
```

Inner Loop:

$10 * \text{mov} + 5 * \text{lea} + 5 * \text{add/inc}$
 $+ 4 * \text{div/mul} + 5 * \text{cmp/br/jmp}$
 $= 29 \text{ instructions}$

Execution time = 43 sec

Optimized Code

```
xorl %r8d, %r8d
xorl %ecx, %ecx
movl %edx, %r9d
cmpl %edx, %r8d
jg .L7
sall $2, %edi
movl %edi, %eax
cltd
idivl %esi
leal 1(%rcx), %edx
movl %eax, %r10d
imull %ecx, %r10d
movl %edx, %ecx
imull %edx, %ecx
leal (%r10, %rcx), %eax
movl %edx, %ecx
addl %eax, %r8d
cmpl %r9d, %edx
jle .L5
.L7: movl %r8d, %eax
ret
```

$4 * \text{mov} + 2 * \text{lea} + 1 * \text{add/inc} +$
 $3 * \text{div/mul} + 2 * \text{cmp/br/jmp}$
 $= 12 \text{ instructions}$

Execution time = 17 sec

Example (input program)

```
int sumcalc(int a, int b, int N)
{
    int i, x, y;
    x = 0;
    y = 0;
    for(i = 0; i <= N; i++) {
        x = x + (4*a/b)*i + (i+1)*(i+1);
        x = x + b*y;
    }
    return x;
}
```

Optimization Example

```
int sumcalc(int a, int b, int N)
{
    int i;
    int x, y;
    x = 0;
    y = 0;
    for(i = 0; i <= N; i++) {
        x = x + (4*a/b)*i + (i+1)*(i+1);
        x = x + b*y;
    }
    return x;
}
```

Example (Output assembly code)

```
sumcalc:  
    pushq  %rbp  
    movq  %rsp, %rbp  
    movl  %edi, -4(%rbp)  
    movl  %esi, -8(%rbp)  
    movl  %edx, -12(%rbp)  
    movl  $0, -20(%rbp)  
    movl  $0, -24(%rbp)  
    movl  $0, -16(%rbp)  
.L2:   movl  -16(%rbp), %eax  
    cmpl  -12(%rbp), %eax  
    jg    .L3  
    movl  -4(%rbp), %eax  
    leal  0(%rax,4), %edx  
    leaq  -8(%rbp), %rax  
    movq  %rax, -40(%rbp)  
    movl  %edx, %eax  
    movq  -40(%rbp), %rcx  
    cltd  
    idivl (%rcx)  
    movl  %eax, -28(%rbp)  
    movl  -28(%rbp), %edx  
    imull -16(%rbp), %edx  
    movl  -16(%rbp), %eax  
    incl  %eax  
    imull %eax, %eax  
    addl  %eax, %edx  
    leaq  -20(%rbp), %rax  
    addl  %edx, (%rax)  
    movl  -8(%rbp), %eax  
    movl  %eax, %edx  
    imull -24(%rbp), %edx  
    leaq  -20(%rbp), %rax  
    addl  %edx, (%rax)  
    leaq  -16(%rbp), %rax  
    incl  (%rax)  
    jmp   .L2  
.L3:   movl  -20(%rbp), %eax  
    leave  
    ret  
  
.size  sumcalc, .-sumcalc  
.section  
.Lframe1:  
    .long  .LECIE1-.LSCIE1  
.LSCIE1:.long  0x0  
    .byte  0x1  
    .string  ""  
    .uleb128 0x1  
    .sleb128 -8  
    .byte  0x10  
    .byte  0xc  
    .uleb128 0x7  
    .uleb128 0x8  
    .byte  0x90  
    .uleb128 0x1  
    .align 8  
.LECIE1:.long  .LEFDE1-.LASFDE1  
    .long  .LASFDE1-.Lframe1  
    .quad  .LFB2  
    .quad  .LFE2-.LFB2  
    .byte  0x4  
    .long  .LCFI0-.LFB2  
    .byte  0xe  
    .uleb128 0x10  
    .byte  0x86  
    .uleb128 0x2  
    .byte  0x4  
    .long  .LCFI1-.LCFI0  
    .byte  0xd  
    .uleb128 0x6  
    .align 8
```

```
pushq    %rbp
movq    %rsp, %rbp
movl    %edi, -4(%rbp)
movl    %esi, -8(%rbp)
movl    %edx, -12(%rbp)
movl    $0, -20(%rbp)
movl    $0, -24(%rbp)
movl    $0, -16(%rbp)
.L2:   movl    -16(%rbp), %eax
cmpl    -12(%rbp), %eax
jg     .L3
movl    -4(%rbp), %eax
leal    0(%rax, 4), %edx
leaq    -8(%rbp), %rax
movq    %rax, -40(%rbp)
movl    %edx, %eax
movq    -40(%rbp), %rcx
cltd
idivl   (%rcx)
movl    %eax, -28(%rbp)
movl    -28(%rbp), %edx
imull   -16(%rbp), %edx
movl    -16(%rbp), %eax
incl    %eax
imull   %eax, %eax
addl    %eax, %edx
leaq    -20(%rbp), %rax
addl    %edx, (%rax)
movl    -8(%rbp), %eax
movl    %eax, %edx
imull   -24(%rbp), %edx
leaq    -20(%rbp), %rax
addl    %edx, (%rax)
leaq    -16(%rbp), %rax
incl    (%rax)
jmp     .L2
.L3:   movl    -20(%rbp), %eax
leave
ret
```

Lets Optimize...

```
int sumcalc(int a, int b, int N)
{
    int i, x, y;
    x = 0;
    y = 0;
    for(i = 0; i <= N; i++) {
        x = x + (4*a/b)*i + (i+1)*(i+1);
        x = x + b*y;
    }
    return x;
}
```

Constant Propagation

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
    x = x + b*y;  
}  
return x;
```

Constant Propagation

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
    x = x + b*y;  
}  
return x;
```

Constant Propagation

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
    x = x + b*0;  
}  
return x;
```

Algebraic Simplification

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
    x = x + b*0;  
}  
return x;
```

Algebraic Simplification

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
    x = x + b*0;  
}  
return x;
```

Algebraic Simplification

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
    x = x;  
}  
return x;
```

Copy Propagation

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
    x = x;  
}  
return x;
```

Copy Propagation

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
    x = x;  
}  
return x;
```

Copy Propagation

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
}  
return x;
```

Common Subexpression Elimination

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
}  
return x;
```

Common Subexpression Elimination

```
int i, x, y;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    x = x + (4*a/b)*i + (i+1)*(i+1);  
}  
return x;
```

Common Subexpression Elimination

```
int i, x, y, t;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + (4*a/b)*i + t*t;  
}  
return x;
```

Dead Code Elimination

```
int i, x, y, t;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + (4*a/b)*i + t*t;  
}  
return x;
```

Dead Code Elimination

```
int i, x, y, t;  
x = 0;  
y = 0;  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + (4*a/b)*i + t*t;  
}  
return x;
```

Dead Code Elimination

```
int i, x, t;  
x = 0;  
  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + (4*a/b)*i + t*t;  
}  
return x;
```

Loop Invariant Code Removal

```
int i, x, t;  
x = 0;  
  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + (4*a/b)*i + t*t;  
}  
return x;
```

Loop Invariant Code Removal

```
int i, x, t;  
x = 0;  
  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + (4*a/b)*i + t*t;  
}  
return x;
```

Loop Invariant Code Removal

```
int i, x, t, u;  
x = 0;  
u = (4*a/b);  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + u*i + t*t;  
}  
return x;
```

Strength Reduction

```
int i, x, t, u;  
x = 0;  
u = (4*a/b);  
  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + u*i + t*t;  
}  
return x;
```

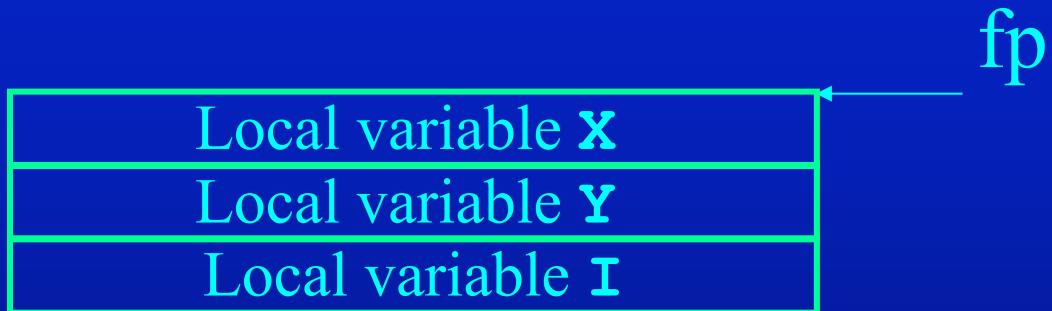
Strength Reduction

```
int i, x, t, u;  
x = 0;  
u = (4*a/b);  
  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + u*i + t*t;  
}  
return x;
```

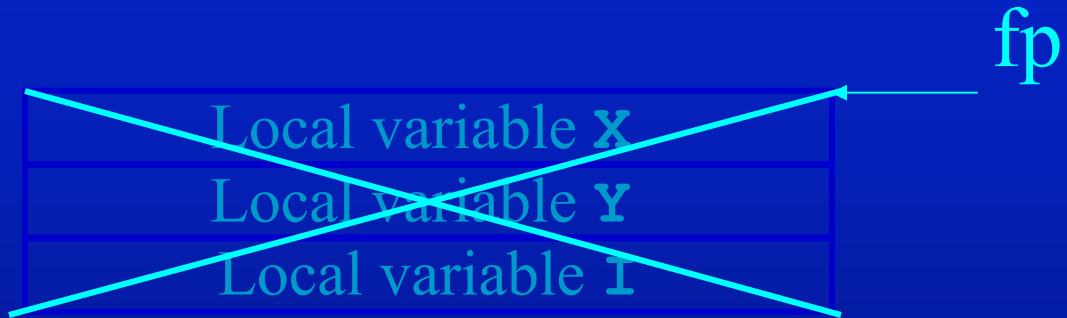
Strength Reduction

```
int i, x, t, u, v;  
x = 0;  
u = ((a<<2)/b);  
v = 0;  
for(i = 0; i <= N; i++) {  
    t = i+1;  
    x = x + v + t*t;  
    v = v + u;  
}  
return x;
```

Register Allocation



Register Allocation



```
$r8d = x  
$r9d = t  
$r10d = u  
$ebx = v  
$ecx = i
```

Optimized Example

```
int sumcalc(int a, int b, int N)
{
    int i, x, t, u, v;
    x = 0;
    u = ((a<<2)/b);
    v = 0;
    for(i = 0; i <= N; i++) {
        t = i+1;
        x = x + v + t*t;
        v = v + u;
    }
    return x;
}
```

Unoptimized Code

```
pushq %rbp
movq %rsp, %rbp
movl %edi, -4(%rbp)
movl %esi, -8(%rbp)
movl %edx, -12(%rbp)
movl $0, -20(%rbp)
movl $0, -24(%rbp)
movl $0, -16(%rbp)
.L2: movl -16(%rbp), %eax
cmpl -12(%rbp), %eax
jg .L3
movl -4(%rbp), %eax
leal 0(%rax, 4), %edx
leaq -8(%rbp), %rax
movq %rax, -40(%rbp)
movl %edx, %eax
movq -40(%rbp), %rcx
cltd
idivl (%rcx)
movl %eax, -28(%rbp)
movl -28(%rbp), %edx
imull -16(%rbp), %edx
movl -16(%rbp), %eax
incl %eax
imull %eax, %eax
addl %eax, %edx
leaq -20(%rbp), %rax
addl %edx, (%rax)
movl -8(%rbp), %eax
movl %eax, %edx
imull -24(%rbp), %edx
leaq -20(%rbp), %rax
addl %edx, (%rax)
leaq -16(%rbp), %rax
incl (%rax)
jmp .L2
.L3: movl -20(%rbp), %eax
leave
ret
```

Inner Loop:

$10 * \text{mov} + 5 * \text{lea} + 5 * \text{add/inc}$
 $+ 4 * \text{div/mul} + 5 * \text{cmp/br/jmp}$
 $= 29 \text{ instructions}$

Execution time = 43 sec

Optimized Code

```
xorl %r8d, %r8d
xorl %ecx, %ecx
movl %edx, %r9d
cmpl %edx, %r8d
jg .L7
sall $2, %edi
movl %edi, %eax
cltd
idivl %esi
leal 1(%rcx), %edx
movl %eax, %r10d
imull %ecx, %r10d
movl %edx, %ecx
imull %edx, %ecx
leal (%r10, %rcx), %eax
movl %edx, %ecx
addl %eax, %r8d
cmpl %r9d, %edx
jle .L5
.L7: movl %r8d, %eax
ret
```

$4 * \text{mov} + 2 * \text{lea} + 1 * \text{add/inc} +$
 $3 * \text{div/mul} + 2 * \text{cmp/br/jmp}$
 $= 12 \text{ instructions}$

Execution time = 17 sec