

# 6.110 Computer Language Engineering

## Recitation 7: Phase 4 infosession

April 5, 2024

**Weekly updates ←**

Phase 4 info

# Wrapping up phase 3...

- **Project phase 3 is due today 11:59PM!!!**
  - **This includes the report!**
  - Remember to add your teammates to the submission!
  - If you need last-minute help, please come to OH today from 2-7pm.

# New releases

- Project phase 4 has been released, due **Friday, April 19**
- Phase 3 LLM questionnaire due **Monday, April 8**
  - Counts as extra credit towards participation grade
- Miniquiz (will be posted soon) and Weekly Check-in are due **Thursday, April 11**
  - Reminder: these are graded on completion – please submit!!

# Lecture forecast... **Week N+1**

<b>Mon 4/8</b>	<b>Tue 4/9</b>	<b>Wed 4/10</b>	<b>Thu 4/11</b>	<b>Fri 4/12</b>
<b>Lectures</b> Foundations of Dataflow ( <i>will take approximately 3 days</i> )				<b>No recitation</b> (CPW)
<b>Due:</b> Phase 3 LLM questionnaire		<b>Re-lecture</b> Optimizations	<b>Due:</b> Mini-quiz, weekly check-in	

# Lecture forecast... **Week N+2**

<b>Mon 4/15</b>	<b>Tue 4/16</b>	<b>Wed 4/17</b>	<b>Thu 4/18</b>	<b>Fri 4/19</b>
<b>Holiday</b> Patriots' Day	<b>No lecture</b>	<b>Guest Lecture</b> Yaron Minsky (Jane Street)	<b>No lecture</b>	<b>Recitation</b> Phase 5 infosession
		<b>Re-lecture</b> Foundations of dataflow	<b>Due:</b> Mini-quiz, weekly check-in	<b>Due: Project phase 4</b>

Weekly updates

**Phase 4 info ←**

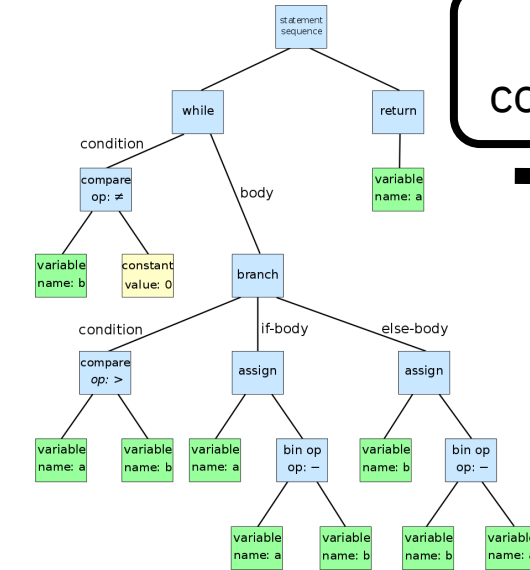
# Project overview

```
import printf;  
void main() {  
  ...  
}
```

**Decaf source file**

**Phase 1.** Does it have the right structure? (syntax)

**Phase 2.** Does it make sense? (semantics)



**Internal representation**

**Phase 3**  
code generation

```
push %rbp  
mov  %rsp, %rbp  
...
```

**x86-64 assembly**



So we have a working compiler now...  
**what next?**

\* Or by the end of today

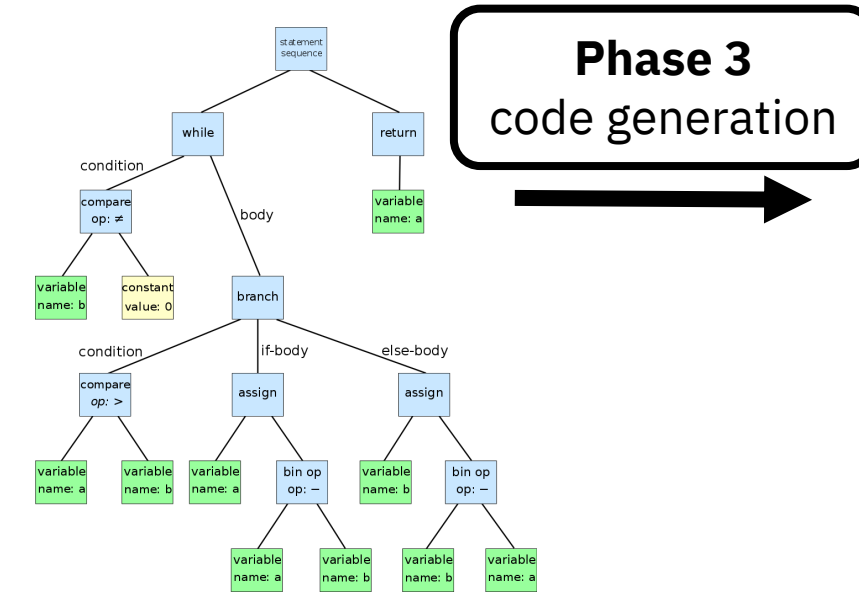
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**Internal representation**

**Phase 3**  
code generation

```
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mov  %rsp, %rbp  
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```

**Optimized  
x86-64 assembly**

**Phase 4.** What can we learn about the program? (dataflow analysis)

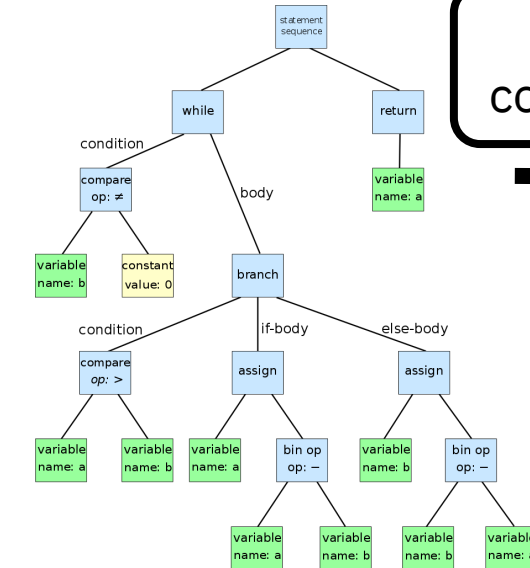
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**Decaf source file**

**Phase 1.** Does it have the right structure? (syntax)

**Phase 2.** Does it make sense? (semantics)



**Internal representation**

**Phase 3**  
code generation

```
push %rbp  
mov  %rsp, %rbp  
...
```

**Even more optimized  
x86-64 assembly**

**Phase 5.** How can we make the output code faster?

**Phase 4.** What can we learn about the program? (dataflow analysis)

From now on, the project becomes more open-ended.

We'll require some specific optimizations, but other than that you are free to implement whatever your heart desire.

At the end of phase 5, there will be a **compiler derby** to find which team's compiler produces the fastest code!

# **Logistics and requirements**

# Phase 4 overview

- **Required:** implement **at least one of the following global dataflow optimizations**
  - Copy propagation
  - Common subexpression elimination
  - Dead code elimination
- Optimization should **at least work on statements involving local (non-array) variables**

# Dataflow analysis: overview

- A form of **program analysis**: compile-time reasoning about program behavior
- Store **some information** we've learned about the program at each program point (CFG node)
- At each node, need to update information based on content of the node (**“transfer function”**), and propagate information to successor nodes (*or predecessors for backwards analyses*)
- At merge points, need to **combine** information somehow
- Iterate until we reach a fixed point
- *More of this formalization in next week's lectures!*

# Copy propagation

- Propagate copies (assignments like  $a \leftarrow b$ )
- Based on **Reaching definitions analysis**: which definitions of each variable reaches each program point\*

$\begin{array}{l} a \leftarrow b \\ c \leftarrow a + 1 \end{array}$	$\begin{array}{l} a \leftarrow b \\ c \leftarrow b + 1 \end{array}$
Before	After



# Copy propagation

- Be careful about this!

a ← b

b ← c

d ← a



a ← b

b ← c

d ← b ???

- One way to avoid: just keep track of which variables are copies of each other instead of using reaching definitions

# Dead code elimination

- Remove code that computes variables that are not used
- Based on **Liveness analysis**: which variables are “live” (has a use afterwards)

$\begin{aligned} a &\leftarrow x + y \\ x &\leftarrow a + b \end{aligned}$ <p>(a is global, x is local)</p>	$a \leftarrow x + y$
Before	After

# Common subexpression elimination

- Only compute an expression once
- Based on **Available expressions analysis**: which expressions defined earlier are still valid (operands not modified)

$\begin{array}{l} a \leftarrow x + y \\ b \leftarrow x + y \\ \quad x \leftarrow a \\ c \leftarrow x + y \end{array}$	$\begin{array}{l} t1 \leftarrow x + y \\ \quad a \leftarrow t1 \\ \quad b \leftarrow t1 \\ \quad x \leftarrow a \\ c \leftarrow x + y \end{array}$
Before	After

# Summary

Optimization	Analysis
Copy propagation	Reaching definitions* <small>*be careful</small>
Common subexpression elimination	Available expressions
Dead code elimination	Liveness

# Summary

	Reaching Definitions	Live Variables	Available Expressions
Domain	Sets of definitions	Sets of variables	Sets of expressions
Direction	Forwards	Backwards	Forwards
Transfer function	$gen_B \cup (x - kill_B)$	$use_B \cup (x - def_B)$	$e\_gen_B \cup (x - e\_kill_B)$
Boundary	$OUT[ENTRY] = \emptyset$	$IN[EXIT] = \emptyset$	$OUT[ENTRY] = \emptyset$
Meet ( $\wedge$ )	$\cup$	$\cup$	$\cap$
Equations	$OUT[B] = f_B(IN[B])$ $IN[B] =$ $\bigwedge_{P, pred(B)} OUT[P]$	$IN[B] = f_B(OUT[B])$ $OUT[B] =$ $\bigwedge_{S, succ(B)} IN[S]$	$OUT[B] = f_B(IN[B])$ $IN[B] =$ $\bigwedge_{P, pred(B)} OUT[P]$
Initialize	$OUT[B] = \emptyset$	$IN[B] = \emptyset$	$OUT[B] = U$

Figure 9.21: Summary of three data-flow problems

# Phase 4 overview (cont'd)

- **Optional:** extend optimizations to global variables and array variables
- **Optional:** other optimizations (more info in handout)
  - Constant propagation and folding
  - Loop-invariant code motion
  - Unreachable code elimination
  - Algebraic simplification (*not dataflow*)
  - ...

# Submission and grading

- Phase 4 is worth **10%** of the overall grade, due **Friday, April 19.**
- Two items to be submitted on Gradescope
  - Design document (8%)
    - Overall dataflow framework (3%)
    - Details of implemented dataflow optimizations (4%)
    - Extras, difficulties, and contributions (1%)
  - Code submission, autograded on correctness only (2%)
    - No private test cases
    - Output code should be correct with and without optimizations

# Specifications

- Your compiler should be **correct** with or without optimizations
- When running  
`./run.sh <filename> -t assembly`  
on a valid input file:
  - Outputs x86-64 assembly code to the output file (or stdout if `-o` is not specified)
- We'll assemble using  
`gcc -O0 -no-pie output.s -o output.exe`



# CLI for optimizations

- **-O cse** turns on common subexpression elimination only
- **-O dce** turns on dead code elimination only
- **-O cp,cse** turns on copy propagation and common subexpression elimination only
- **-O all** **turns on all optimizations** (we'll run the autograder with this option)
- **-O all,-cse** turns on all optimizations except common subexpression elimination

# Design document

- Explains technical details
- Includes the following sections:
  1. Design (*including general dataflow framework and specific details for each implemented optimization*)
  2. Extras
  3. Difficulties
  4. Contribution
- If you used LLMs, also describe how you used them and provide chat logs

# 1. Design

- Overview of your design, including design choices you made and design alternatives you considered.
- This section should help us understand your code
- In particular, please include:
  - Your general framework for dataflow optimizations (worth **3%**)
  - Details of each dataflow optimization you implemented (worth **4%**, more info on next slide)

# 1. Design — details

- For each dataflow optimization you implemented, please include:
  - the scope of the optimization (did you take into account global variables and/or array variables?)
  - the dataflow equations you used
  - a sample test case, with generated code before and after, included under [doc/phase4-code/](#) in your repository
  - a brief explanation of how your dataflow optimization worked

# Other sections (worth **1%**)

## 2. Extras:

- Any clarifications, assumptions, or additions you made
- Any interesting debugging techniques, build scripts
- Anything cool you'd like to share!

## 3. Difficulties:

- List of known problems with your project, and as much as you know about the cause
- Any issues from phase 3 that you fixed

## 4. Contributions: A brief description of how your group divided the work

**Words of advice**

- **Start simple!**

- Start with very simple test cases so that you understand what's happening
- Start with local non-array variables only, and only add global variables / array variables after you can get the analysis to work on local variables

- **Keep things general**
  - Various dataflow analyses can all be written in terms of a transfer function and a meet function
  - Consider making a parametrized dataflow framework
  - *Next week's lecture will cover this formalization*

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Figure 9.21: Summary of three data-flow problems



- **Consider using single-statement blocks**
  - More time/memory-consuming but who cares
  - No need to propagate information inside a basic block
  - One tricky thing: Need to be able to add/remove nodes/merge points/join points.

- **Use array of nodes, not pointer-and-objects**
  - Key: Need to be able to remove/add statements
  - Especially relevant if you don't use basic blocks
  - You will need adjacency list and reverse adj. list